

Quanterall HQ Varna, Bulgaria 2019

Namdak Tonpa

The Languages

Groupoid Infinity

About Speaker

- PhD student, 3-rd year of education (<https://cubical.systems>)
- Author of 8 programming languages and 2 runtime cores
- But more famous for N2O framework (<https://n2o.dev>)
- Love to create programming languages and talk about them
- Know how to convert open source to money
- Aware of all operating systems/programming languages (~100/~1000)

Github Organizations

- GROUPOID – The Language of Space
- SYNRC – Application Layer Formal Specification and Implementations
- VOXOZ – Virtual Machines and Network Infrastructure

Talk Structure

The Languages

I. Languages

- Main Contributions
- Industrial Compilers
- Fast Interpreters
- Formal Verification

II. Processing

- History
- Workflow Languages
- Financial Languages
- Contract Languages

Main Contributions

- John McCarthy [LISP]
 - Robin Milner [ML, Pi Calculus, HOL]
 - Simon Peyton Jones [Haskell]
 - Xavier Leroy [OCaml]
 - Niko Matsakis [Rust] Linear Types
 - Joe Armstrong [Erlang] ... and many others
-
- Nicolaas Govert de Bruijn [AUTOMATH]
 - Thierry Coquand [Coq]
 - Ulf Norell [Agda]
 - Leonardo de Moura [Lean] ... and not so many

Industry

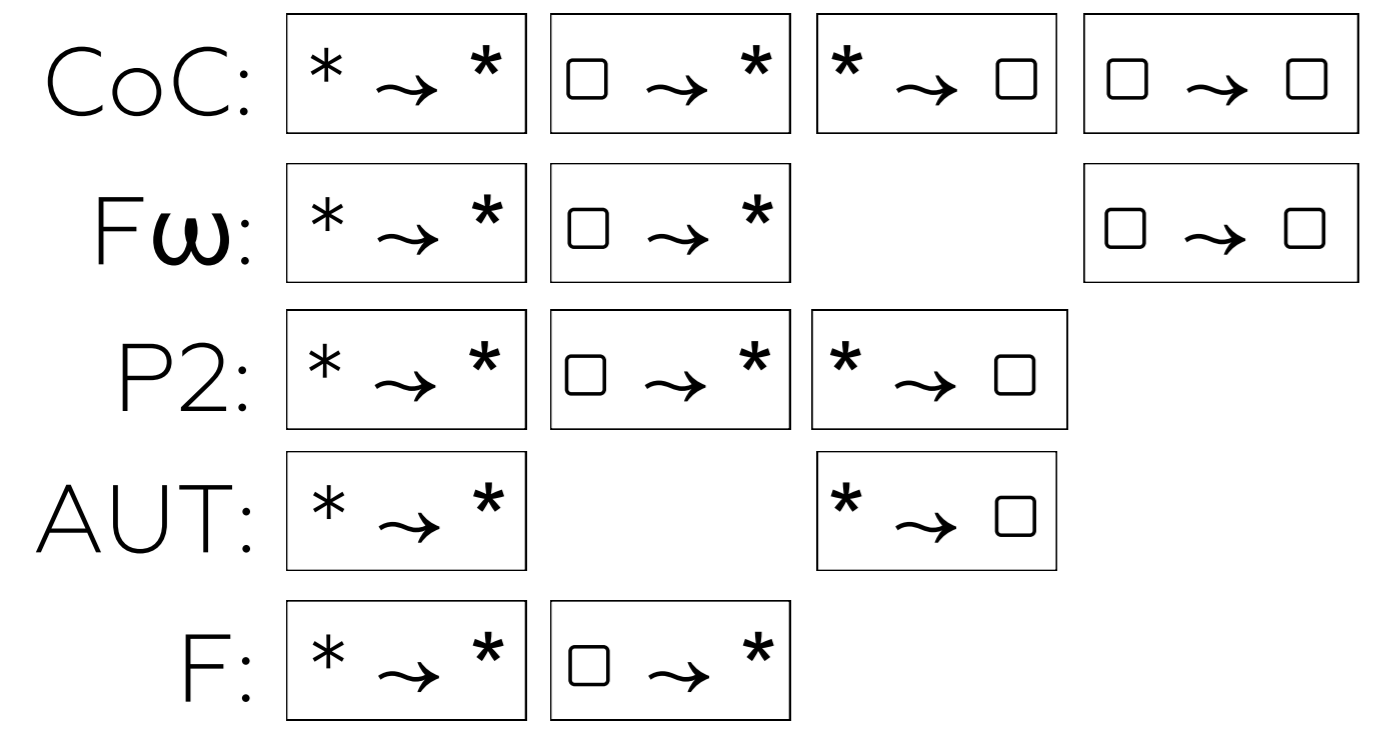
- V8, WebAssembly (any)
- LuaJIT (nginx)
- JVM (Oracle)
- CLR (MS) ... and other JITs
- IR/MIR/LLVM (clang, rust)
- OCaml/GHC
- SPIRAL

MOTTO 1: If you have compact language that fits L1 cache along with its interpreter, then you don't need JIT! However you still need vectorization.

MOTTO 2: At enterprise scale you still need types or ULC targeted extraction.

λ -Calculi

in Extended Lambda Cube



System F ω :
 Haskell, Scala, 1ML
 Almost CoC
 No Types
 On Values

System F:
 ML, Miranda,
 OCaml

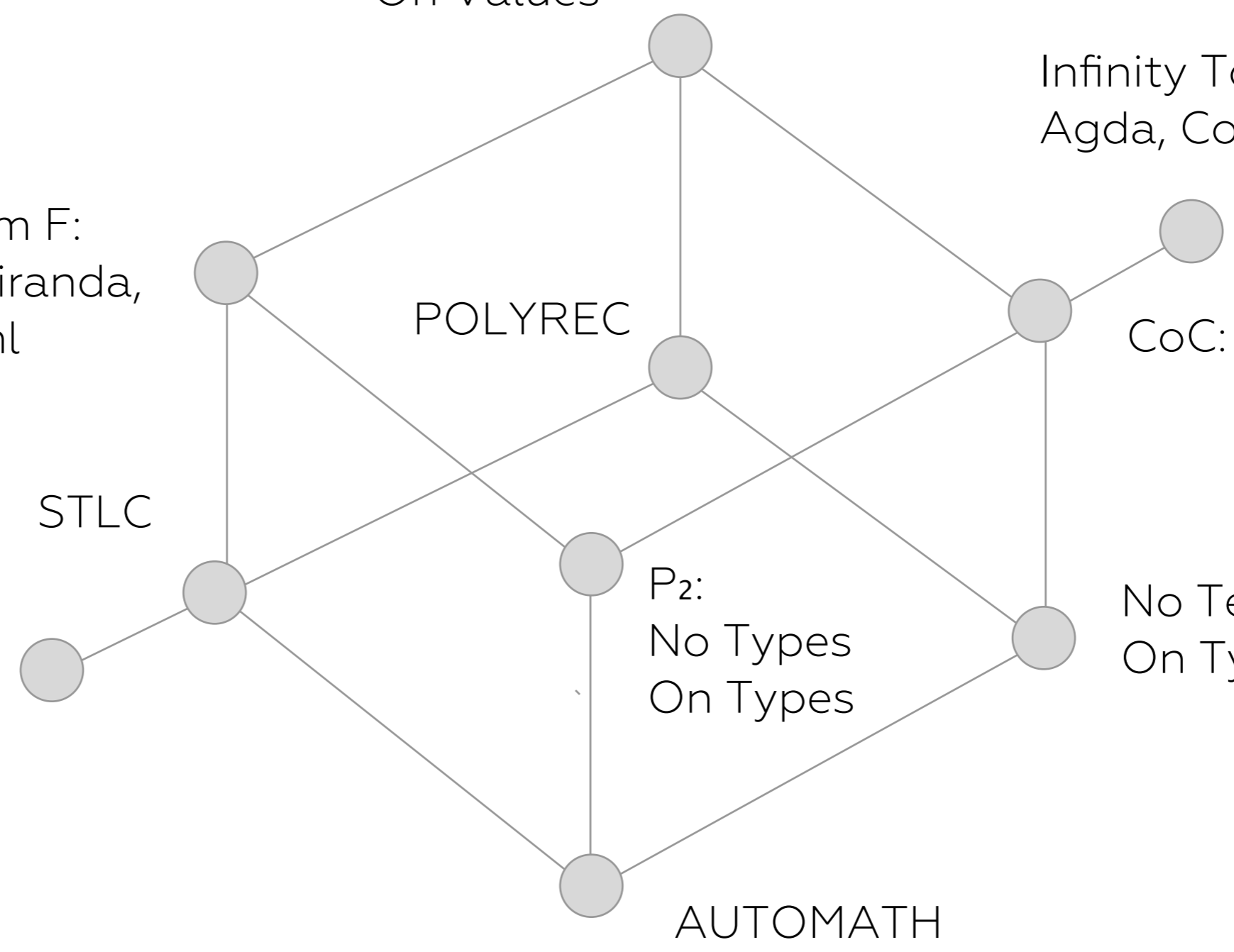
Infinity Topoi
 Agda, Coq, Lean, Om

CoC: Morte, Henk

Untyped SLC:
 Erlang, LISP,
 JavaScript

P₂:
 No Types
 On Types

No Terms
 On Types



Formal Verification

Mathematical Formal Software Verification unveils the inner structure of phenomena and avoid wide range of errors.

- 1) Mars Climate Orbiter (1998), conversion inch/met — \$80M;
 - 2) Ariane Rocket (1996), downcast from 64 to 16 bit — \$500M;
 - 3) FPU DIV Error Pentium (1994) — \$300M;
 - 4) Business Contract Error EVM — \$50M;
 - 5) Error in SSL (heartbleed) — \$400M.
-
- 1) IEEE Std 1012-2016 — V&V Software verification and validation (4 layers)
 - 2) ESA PSS-05-10 1-1 1995 — Guide to software verification and validation;
 - 3) ISO/IEC 13568:2002 — Z formal specification notation.

Attempts to Fix C/C++

Expensive and long way of doing things...

- Coq: VST, DeepSpec
- Haskell, HOL: L4
- Even Manual Proofs!!!

Deep Embedding

... seems a better way exist – direct certified extraction with no intermediate proofs!

- Coq: The best macroassembler
- Coq.io – OCaml/Lwt bindings
- Agda x86
- Clash, Lambda to VDHL/Verilog

History of Processing Languages

- EMAIL: FSM
- Event-Condition-Action Reactive Rule Engines
- Expert Systems: RETE Engine, Prolog
- Workflow Standards of the past: XPD, BPML, OpenWFE, WWF and jBPM
- Workflow Standard After 2008: BPMN
- Trading: TpML, Fix, business contract routers, cross-exchanges, arbitrage
- Business Contracts Virtual Machines: EVM, Script VM, aebytecode
- Business Contract Languages: Sophia, Solidity, Plutus
- MLTT Frameworks: Dhall
- Interaction Networks Evaluators: Formality, Moonad
- Stream Processing: Oz, Erlang, np/ling, Futhark

What is the Language?

Prerequisites for bootstrapping are algebraic data types: struct (*) and union (+) from C/C++

Logic Core:

```
data pts = star (n: nat)
         | var (l: nat)
         | pi (l: nat) (d c: pts)
         | lambda (l: nat) (d c: pts)
         | app (f a: pts)
```

Runtime Core:

```
data ulc = var (l: nat)
         | lambda (l: nat) (d c: ulc)
         | app (f a: ulc)
```

Is that enough?

No, we need Inductive Types!

Inductive Core:

```
data tele (A: U) = emp | tel (n: name) (b: A) (t: tele A)
data branch (A: U) = br (n: name) (args: list name) (term: A)
data label (A: U) = lab (n: name) (t: tele A)
data ind = data_ (n: name) (t: tele lang) (labels: list (label lang))
          | case (n: name) (t: lang) (branches: list (branch lang))
          | ctor (n: name) (args: list lang)
```

IO

And we need Effects to access to business rules!

IO Core:

```
data IO (A: U)
  = getLine (_: String → IO)
  | putLine (_: String)
  | pure (_: A)
```

Secure Storage:

```
data KV (A: U)
  = get (_: String → IO)
  | put (_: String)
  | sign (_: String → IO)
  | verify (_: String → IO)
  | pure (_: A)
```

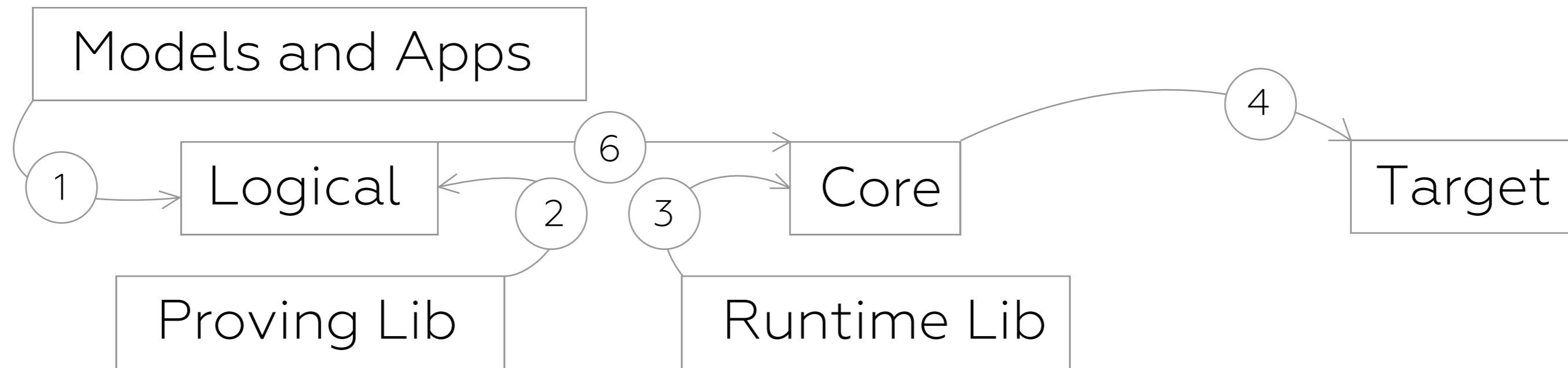
Infinity IO

What about Infinitary IO?

```
data IOI.F (A State: U) = getLine (_: String → State)
                          | putLine (_: String) (_: State)
                          | pure (_: A)
```

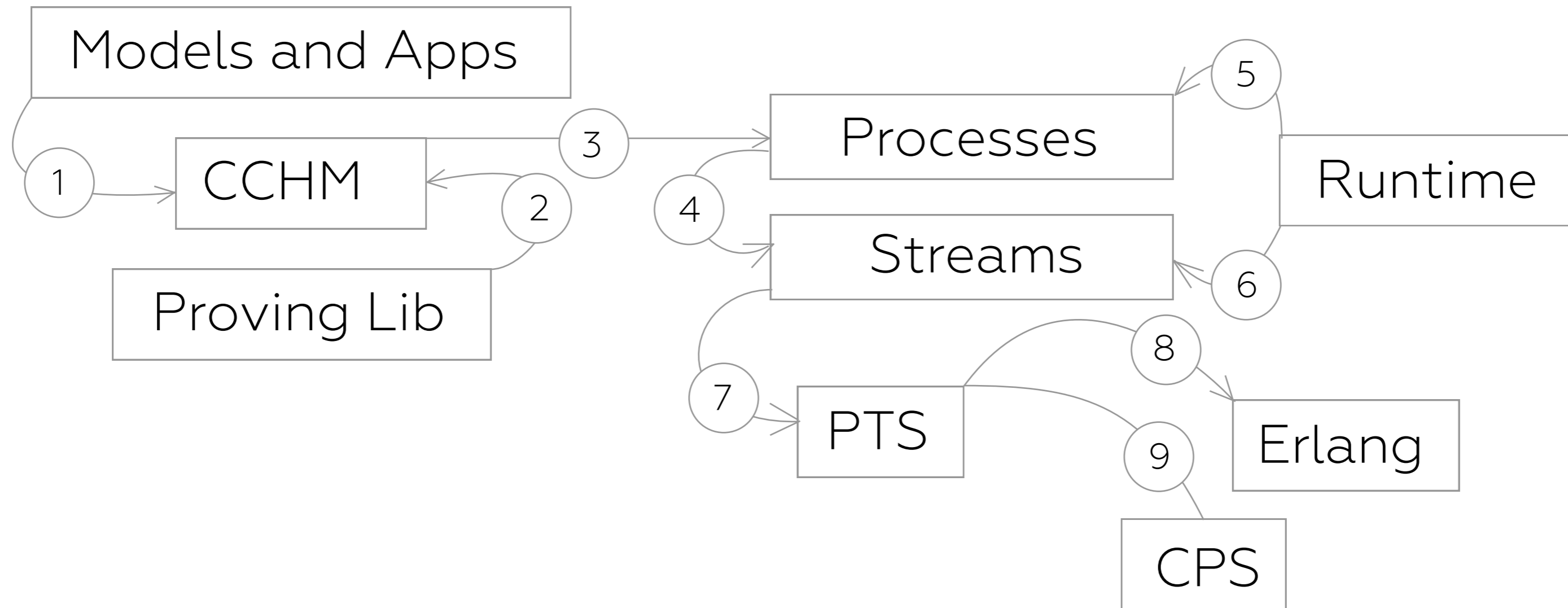
```
data IOI (A State U) = intro (_: State). (_: State -> IOI.F A State)
```


Verification Process #1



1. Model Specification
2. Model Checking
3. Runtime Linkage
4. Target Machine Code Extraction

Verification Process #2



Research Subject

Classification of Languages use in
Specification, Formalization and Verification process

- 1) Specification Languages (Z, UML, MLTT);
- 2) Model Checkers (TLA+, Twelf, Dedukti, Z3);
- 3) General Purpose Languages (Haskell, OCaml, Erlang, Scala, LISP);
- 4) Theorem Provers (Agda, Coq, HOL, ACL2);
- 5) Unified Execution Environments (HaLVM, LING, Mirage);
- 6) Contract Machines and Languages (EVM, Script VM, Sophia, Plutus)
- 7) Workflow Languages (BPMN)
- 9) Exchange Trading Languages (TpML)

Plutus Review

IOHK Certified Language for Haskell Embedding and Development

- 1) Certification and Formalization (Agda): NbE, Extraction
- 2) Plutus IR (Lisp): Intermediate Language, Fix, No Pattern Match Compiler
- 3) Plutus Core: CEK, L machines
- 4) Scott Encoding of Data Types
- 5) Marlowe: Business Contracts (Alexander Nemish)
- 6) Plutus TxCompiler: Haskell Code to Plutus (getPlc)

Plutus IR AST

```
data Term tynome name a
  = Let a Recursivity [Binding tynome name a] (Term tynome name a)
  | Var a (name a)
  | TyAbs a (tynome a) (Kind a) (Term tynome name a)
  | LamAbs a (name a) (Type tynome a) (Term tynome name a)
  | Apply a (Term tynome name a) (Term tynome name a)
  | Constant a (PLC.Constant a)
  | Builtin a (PLC.Builtin a)
  | TyInst a (Term tynome name a) (Type tynome a)
  | Error a (Type tynome a)
  | IWrap a (Type tynome a) (Type tynome a) (Term tynome name a)
  | Unwrap a (Term tynome name a)
```

Plutus IR Sample

IOHK Certified Language for Haskell Embedding and Development

```
(lam pubkey (con bytestring)
 (lam signed (con bytestring)
  [ { (abs a (type) (lam b (all a (type) (fun a (fun a a))))
      (lam t (fun (all a (type) (fun a a)) a) (lam f (fun (all a (type) (fun a a)) a)
        [ [ { b (fun (all a (type) (fun a a)) a } t f ] (abs a (type) (lam x a x)) ] ) ) ) )
      (all a (type) (fun a a)) } [ (builtin verifySignature) signed txhash pubkey ] (lam
u (all a (type) (fun a a)) (abs a (type) (lam x a x)) )
(lam u (all a (type) (fun a a)) (error (all a (type) (fun a a)))) ] ) )
```

Pure Core/CoC/Morte/Om

Theoretical Mimimum Scholarship Language Development
Toy Dependently Typed Language for Typechecking and Extraction

- CoC, Morte, Om (Pure Core)
- Further Evolution (Inductive Types): Dhall, Formality

Specially created for Erlang deployment!

Real Monads Extracted from CoC to Erlang bytecode!

> 'Monad':

```
'[<=<]'/0    '[=<<]'/0    '[>=>]'/0    '[>>=]'/0    forM/0  
forM_/0      join/0      mapM/0      mapM_/0      module_info/0  
module_info/1 replicateM/0 replicateM_/0 sequence/0 sequence_/0
```

Formality

Interaction Networks based Evaluator (Run-time Fusion)
GPU Backend, Rust Implementation
Faster than GHC

```
id(10000000000(List<Bool>, map(Bool, Bool, not), list))
```

Flips every bit in a list of 100 bits, a billion times. It prints the correct output in 0.03s. You could increase that to beyond the number of stars in the universe, and it'd still output the correct result, instantly.

<https://github.com/moonad/whitepaper>